

GENERAL EQUILIBRIUM - HW #1

1. **Gainers and Losers in Equilibrium**

Note that preferences are identical and homothetic, so we may use a representative agent to find PTE's.

- (a) In each region, the PTE allocations are symmetric:

$$\omega_{A_1} = \omega_{A_2} = (2, 1)$$

$$\omega_{B_1} = \omega_{B_2} = (1, 2)$$

To see this in region A , substitute $y = 2z$. This yields symmetric preferences and endowments, $u = xz/2$ and $\omega = (4, 4)$, so the RA sets equal prices: $(p_x, p_z) = (1, 1)$. Equivalently,

$$\mathbf{p} = (p_x, p_y) = (1, 2)$$

With these prices, A_1 and A_2 have equal wealth, so each consumes the half the aggregate endowment, obtaining $u = 2$. The analysis is identical for region B, but with $\mathbf{p} = (2, 1)$.

- (b) Aggregate endowment is symmetric: $\omega = (6, 6)$, so $\mathbf{p} = (1, 1)$. Type 1 agents have twice as much wealth as Type 2 agents, so they consume twice as much of the aggregate endowment:

$$\text{Type 1: } c = (2, 2); u = 4$$

$$\text{Type 2: } c = (1, 1); u = 1$$

- (c-d) Any allocation that gives 1/4 of the total wealth to each agent results in a PTE with each agent consuming (1.5, 1.5) and utility $u = 2.25$.
- (e) The allocation in (b) is Pareto Efficient, as is the allocation in (d). Free trade does not guarantee a "fair" equilibrium, nor one that maximizes aggregate utility.

2. Replica Invariance

- (a) Any utility $u_1 + u_2 \leq 100$ is feasible. From the given endowment, there are no mutually advantageous trades.
- (b) In \mathcal{E}^2 a more efficient allocation is possible:

Type 1: $c = (50, 50)$

Type 2: $c_1 = (100, 0), c_2 = (0, 100)$

The Type 2 agents now have increased utility, $u = 100$, and the Type 1 still have $u = 50$ as in the original endowment.

- (c) In a PTE, each agent maximizes utility subject to his wealth constraint. If there is a PTE in \mathcal{E} , then the same prices must yield a PTE in \mathcal{E}^2 .

As we'll show in (e), the allocation in (b) is a PTE with $\mathbf{p} = 1$. Any PTE in \mathcal{E} must therefore have the same prices. We also know that a PTE in \mathcal{E} must allocation $(50,50)$ for each agent, as this endowment is already Pareto Optimal.

This allocation, however, is not an equilibrium with $\mathbf{p} = 1$, because the Type 2 agents can increase utility by trading. For example, $c = (49, 51)$ yields $u = 51$. It follows that there is no PTE in \mathcal{E} .

- (d) Preferences for the Type 1 agent do not exhibit local non-satiation. For example, at both $(50,50)$ and $(50,51)$, his utility is 50. Preferences for the Type 2 agent are not convex. For example, both $(100,0)$ and $(0,100)$ are strictly preferred to $(50,50)$.
- (e) For \mathcal{F}^k , give $(50,50)$ to each Type 1 agent, and $(100,0)$ to half the Type 2's, $(0,100)$ to the other half. No agent can be made better off without increasing the gross total $x + y$ of his consumption. Thus, we cannot improve upon this allocation without increasing the gross aggregate endowment, $\omega_x + \omega_y$. The allocation is therefore efficient for all k .
- (f) With $\mathbf{p} = 1$, each agent's budget constraint is $x + y = 100$. The Type 1 agents maximize utility at $(50,50)$, while the Type 2 agents maximize at the boundaries $(0,100)$ or $(100,0)$. In the allocation in (e), each agent consumes one of these utility maximizing bundles. The allocation is therefore a PTE.

3. Demand Theory with Quasi-linear Utility

- (a) Denote indirect utility as $V(p, w)$. Let any w, w' be given and denote the corresponding optimal non-money consumption as x, x' respectively. Since x is optimal for w , we have

$$\begin{aligned} U(x, w - px) &\geq U(x', w - px') \\ &= v(x') + w - px' \\ &= U(x', w' - px') + (w - w') \end{aligned}$$

Comparing the first and last lines, we have

$$V(p, w) - w \geq V(p, w') - w',$$

but our choice of w and w' was arbitrary, so the reverse inequality also holds. Thus, for all w, w' , we have

$$V(p, w) - w = V(p, w') - w',$$

So $\alpha(p) = V(p, w) - w$ is independent of w , as desired.

- (b) Denote the expenditure function as $C(p, U)$. From duality theory, we know

$$U = V(p, C(p, U))$$

Applying the form from (a), this becomes

$$U = \alpha(p) + C(p, U),$$

so expenditure is simply

$$C(p, U) = -\alpha(p) - U$$

- (c) Given any w, w' , we show that we achieve the same utility with either x or x' :

$$\begin{aligned} U(x', w - px') &= v(x') + w - px' \\ &= v(x') + w' - px' + (w - w') \\ &= U(x', w') + w - w' \\ &= V(p, w') - w' + w \\ &= V(p, w) - w + w \quad (\text{from (a)}) \\ &= V(p, w) \\ &= U(x, w - px) \end{aligned}$$

Thus, x' is optimal for wealth w , so optimal non-money consumption is independent of wealth, as desired.

(d) In equilibrium, we have

$$\partial v / \partial x_i = p_i$$

By the envelope theorem, differentiating with respect to p_i gives

$$\frac{\partial^2 v}{\partial x_i^2} \frac{\partial x_i}{\partial p_i} = 1$$

Assuming $v(x)$ is concave, its second differential is negative, so $\partial x_i / \partial p_i$ is also negative, as is own price elasticity.

4. **Demand/Supply, Inverse Demand/Supply and Indirect Utility/Profit**

(a) (1 \Rightarrow 2): Suppose $(z, m) \in D(p)$. Then $pz + m = 0$ and

$$v(z) + m \geq v(z') + m'$$

for all (m', z') with $pz' + m' = 0$. So we have

$$v(z) - pz \geq v(z') - pz'$$

for all $z' \in Z$. By definition, $p \in \partial v(z)$.

(2 \Rightarrow 3): If $p \in \partial v(z)$, then for all $z' \in Z$, we have

$$v(z) - pz \geq v(z') - pz'.$$

That is,

$$\begin{aligned} v(z) - pz &= \max\{v(z') - pz'\} \\ &= \max\{v(z') + m' \mid pz' + m' = 0\} \end{aligned}$$

That is, $v(z) - pz$ is the maximum utility feasible with zero total wealth. This is the definition of the indirect utility function $v^*(p)$.

(3 \Rightarrow 1): By definition, $v^*(p)$ is the maximum utility achievable with zero net wealth. That is,

$$v^*(p) = \max\{v(z') + m' \mid pz' + m' = 0\}.$$

Since $v(z) - pz = v^*(p)$, this implies $(z, -pz) \in D(p)$.

(b) The revenue from selling z is $r = -pz$, which we previously denoted m . The firm's profit is therefore

$$\pi(p, z) = -c(z) - pz = v(z) + r$$

In this notation, the conditions above become

$$(1) (z, r) \in \operatorname{argmax}_{z, r} \{v(z) + r \mid pz + r = 0\}$$

$$(2) -p \in \partial c(z)$$

$$(3) \pi^*(p) = -c(z) - pz$$